## Boomerang Connectivity Table: A New Cryptanalysis Tool

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## Differential Cryptanalysis

[Biham-Shamir1990]

- Prepare two input values $P_{1}, P_{2}$ with (usually) small difference $\Delta P=P_{1} \oplus P_{2}$.
- Expecting some output differences $\Delta C=C_{1} \oplus C_{2}$ with a high probability.

Solid methods to evaluate probability are evaluated.

## Differential Cryptanalysis

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## Boomerang Attacks

Proposed by [Wag99] to combine independent two characteristics.

- $E_{0}: \operatorname{Pr}\left[\Delta_{i} \rightarrow \Delta_{o}\right]=p$
- $E_{1}: \operatorname{Pr}\left[\nabla_{i} \rightarrow \nabla_{o}\right]=q$

Two pairs are analyzed. Distinguish probability:

$$
p^{2} q^{2}
$$



## Two Trails in Boomerang Attacks

[Wag99]: Assumed two trails are independent. not always correct

- Dependency can help attackers.
[BDD03]: Middle-round S-box trick
[BK09]: Boomerang switch
Ladder switch / Feistel switch / S-box switch
- Dependency can spoil attacks.
[Mer09]: Incompatible trails


## Ladder Switch



## Ladder Switch



1/1


SB

$E_{0}$ : Columns 3: no active S-box for $E_{0}$ $E_{1}$ : Columns 0: no active S-box for $E_{1}$ © ntt

Feistel Switch / S-box Switch

prob to be a right quartet is $p\left(\operatorname{not} p^{2}\right)$

## Sandwich Attacks [DKS10]

Generalized framework including dependency of two trails:

$$
E=E_{1} \circ E_{m} \circ E_{0}
$$

Distinguish probability is $p^{2} q^{2} r$, with some probability $r$ for $E_{m}$.


## Probability for $E_{m}$

$r=\frac{\#\left\{x \in\{0,1\}^{n} \mid E_{m}^{-1}\left(E_{m}(x) \oplus \nabla_{i}\right) \oplus E_{m}^{-1}\left(E_{m}\left(x \oplus \Delta_{o}\right) \oplus \nabla_{i}\right)=\Delta_{o}\right\}}{2^{n}}$


Probability space is only the size of $E_{m}$, not its square.

Ladder Switch $r=1 \quad$ S-box Switch $r=p$


- $r$ is for a quartet, not for a pair in the standard differential cryptanalysis. How to evaluate it?
- Our focus: $E_{m}$ is a single S-box layer
- a new form to easily evaluate $r$ for S-box

Adv. 1: new switching effect ( $r$ is surprisingly high)
Adv. 2: quantitating the strength of $S$-box against sandwich attack (a new S-box design criterion)

- We reveal several relationships between the standard probability in DDT and $r$.


## DDT: Differential Distribution Table

$$
\begin{aligned}
& \#\left\{x \in\{0,1\}^{n} \mid S(x) \oplus S\left(x \oplus \Delta_{i}\right)=\Delta_{o}\right\} \\
& \#\left\{x \in\{0,1\}^{n} \mid S(x) \oplus S\left(x \oplus \Delta_{i}\right)=\Delta_{o}\right\}
\end{aligned}
$$

## BCT: Boomerang Connectivity Table

$\#\left\{x \in\{0,1\}^{n} \mid S^{-1}\left(S(x) \oplus \nabla_{o}\right) \oplus S^{-1}\left(S\left(x \oplus \Delta_{i}\right) \oplus \nabla_{o}\right)=\Delta_{i}\right\}$


## Observations of BCT (1/3)

Observations of BCT (2/3)

## S-box Switch: $\quad \operatorname{Pr}[\Delta \xrightarrow{S} \nabla]=\boldsymbol{p} " \Rightarrow \boldsymbol{p}^{\boldsymbol{S}}=\boldsymbol{p} "$

Lemma 1 For any choice of $\left(\Delta_{i}, \Delta_{o}\right)$, the value in the BCT is greater than or equal to the one in the DDT.

|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 16 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 16 | 16 | 16 | 16 | 16 | 16 | 16 | 16 |
| 1 | 0 | 0 | 0 | 4 | 0 | 0 | 0 | 4 | 1 | 16 | 0 | 4 | 4 | 0 | 16 | 4 | 4 |
| 2 | 0 | 0 | 0 | 2 | 0 | 4 | 2 | 0 | 2 | 16 | 0 | 0 | 6 | 0 | 4 | 6 | 0 |
| 3 | 0 | 2 | 0 | 2 | 2 | 0 | 4 | 2 | 3 | 16 | 2 | 0 | 6 | 2 | 4 | 4 | 2 |
| 4 | 0 | 0 | 0 | 0 | 0 | 4 | 2 | 2 | 4 | 16 | 0 | 0 | 0 | 0 | 4 | 2 | 2 |
| DDT |  |  |  |  |  |  |  |  | ВСТ |  |  |  |  |  |  |  |  |

S-box switch is the equal case of Lem. 1

Observations of BCT (3/3)

## Values in BCT can be bigger than DDT.

|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 16 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 4 | 0 | 0 | 0 | 4 |
| 2 | 0 | 0 | 0 | 2 | 0 | 4 | 2 | 0 |
| 3 | 0 | 2 | 0 | 2 | 2 | 0 | 4 | 2 |
| 4 | 0 | 0 | 0 | 0 | 0 | 4 | 2 | 2 |


|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 16 | 16 | 16 | 16 | 16 | 16 | 16 | 16 |
| 1 | 16 | 0 | 4 | 4 | 0 | 16 | 4 | 4 |
| 2 | 16 | 0 | 0 | 6 | 0 | 4 | 6 | 0 |
| 3 | 16 | 2 | 0 | 6 | 2 | 4 | 4 | 2 |
| 4 | 16 | 0 | 0 | 0 | 0 | 4 | 2 | 2 |
| BCT |  |  |  |  |  |  |  |  |

Comparison of DDT and BCT for AES S-box

| Value | 256 | 6 | 4 | 2 | 0 |
| :--- | :---: | :---: | :---: | :---: | :---: |
| DDT | 1 | - | 255 | 32130 | 33150 |
| BCT | 511 | 510 | 255 | 31620 | 32640 |

## Generalized Switching Effect

- Focus on $\left(\Delta_{i}, \Delta_{o}\right)$ whose DDT entry is 4.
- 2 pairs satisfying those diff propagation



How can we define $\nabla$ s.t. a quartet is formed?

## Generalized Switching Effect

- 3 ways to define $\nabla$, one is known as $S$-box switch



## Generalized Switching Effect

- 3 ways to define $\nabla$, one is known as $S$-box switch


Lemma 2 For any fixed $\Delta_{i}$, for each entry with ' 4 ' in the DDT, the value of two positions in the BCT will increase by 4.
(O) NTT

## Generalized Switch for 6-uniform DDT

We can make 3 distinct quartets. Each increases the value of BCT in 2 positions.
$x_{5}$


## Applications so far

Related-tweakey boomerang distinguisher on 8round Deoxys-384:

- Prev: $2^{-6}$ (single S-box switch)
- New: $2^{-5.4}$ (single generalized switch)
- 9R and 10R distinguishers are also improved.

Related-tweakey rectangle attacks on SKINNY

- Prev: prob was experimentally evaluated
- New: theoretical analysis of the probability


## Extension to ARX Construction

Similar analysis can be applied to modular addition.
(

Case Study: 3-bit Addition $\left(\Delta_{i}=0\right)$

| DD |  |  |  |  |  |  |  |  |  | BC |  |  |  |  | $\nabla$ o |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |  |  | 0 | 1 | 2 |  | 4 | 5 | 6 |  | 7 |
| 0 | 64 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  | 0 | 64 | 464 | 464 | 46 | 64 | 64 | 464 | 4 | 64 |
| 1 | 0 | 32 | 0 | 16 | 0 | 0 | 0 | 16 |  | 1 | 64 | 0 | ) 32 | 0 | 64 | 0 | 32 |  | 0 |
| 2 | 0 | 0 | 32 | 0 | 0 | 0 | 32 | 0 |  | 2 | 64 | $4 \overline{64}$ | 40 | 0 | 64 | 64 | 40 |  | 0 |
| $\Delta_{i}^{\prime} 3$ | 0 | 16 | 0 | 16 | 0 | 16 | 0 | 16 | $\Delta_{i}^{\prime}$ | 3 | 64 | 40 | 32 | 0 | 64 | 0 | 32 |  | 0 |
| 4 | 0 | 0 | 0 | 0 | 64 | 0 | 0 | 0 |  | 4 | 64 | 464 | 464 | 464 | 64 | 64 | 64 |  | 64 |
| 5 | 0 | 0 | 0 | 16 | 0 | 32 | 0 | 16 |  | 5 | 64 | 40 | 32 | 20 | 64 | 0 | 32 |  | 0 |
| 6 | 0 | 0 | 32 | 0 | 0 | 0 | 32 | 0 |  | 6 | 64 | 464 | 40 | 0 | 64 | 64 | 40 |  | 0 |
|  | 0 | 16 | 0 | 16 | 0 | 16 | 0 | 16 |  | 7 | 64 | 40 | 32 | 20 | 64 | 0 | 32 |  | 0 |

- BCT < DDT (S-box switch does not work)
- MSB switch


## Concluding Remarks

BCT: precomp table of $r$ in the sandwich attack Adv. 1: new switching effect ( $r$ is surprisingly high) Adv. 2: quantitating the strength of S-box against sandwich attack (S-box design criteria)

Problems to investigate

- improving previous boomerang attacks
- extending $E_{m}$ (more than single S-layer)
- comprehensive study for modular addition

